



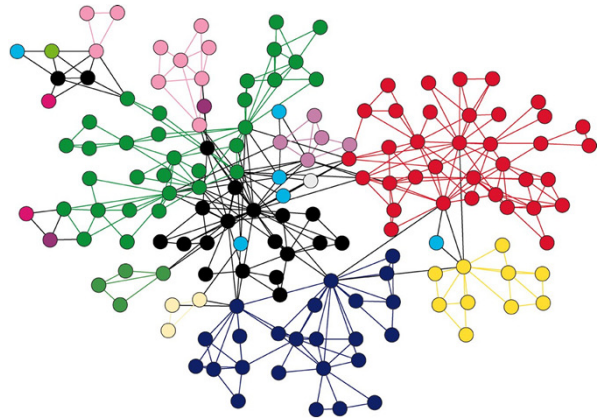
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Data Structure

Lecture 6: The Graph

Asst. Instructor

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Graph

- *A Graph is a non-linear data structure consisting of nodes and **edges**. The nodes are sometimes also referred to as **vertices** and the edges are lines or arcs that connect any two nodes in the graph.*
- *Graphs provide the ultimate in data structure flexibility. Its can model both real-world systems, so they are used in hundreds of applications.*
- *Since they are powerful abstractions, graphs can be very important in modeling data. In fact, Graphs have applications in a host of different domains, including:*
 1. **Social network graphs:** *to tweet or not to tweet. Graphs that represent who knows whom, who communicates with whom, who influences whom or other relationships in social structures.*

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Graph

2. **Transportation networks.** In road networks **vertices** are intersections and **edges** are the road segments between them, and for public transportation networks **vertices** are stops and **edges** are the links between them. Such networks are used by many map programs such as **Google maps** and **GPS** maps.
3. **Utility graphs.** The power grid, the Internet, and the water network are all examples of graphs where **vertices** represent connection points, and **edges** the wires or pipes between them.
4. **Document link graphs.** The best known example is the link graph of the web, where each web page is a vertex, and each hyperlink a directed edge.

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Graph: Formal Def.

- **Graph:** A graph consists of a set of vertices V and a set of edges E , The set of edges describes relationships among the vertices such that each edge in E is a connection between a pair of vertices in V .
- A graph G is defined as follows:

$$G = (V, E)$$

$V(G)$: a finite, nonempty set of vertices.

$E(G)$: a set of edges (pairs of vertices).

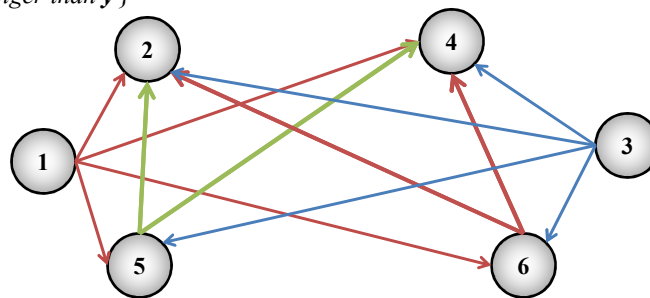
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Graph: Example_1

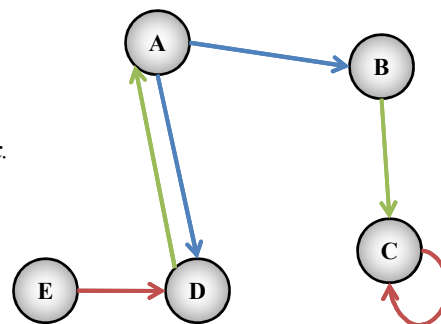
- A “Real-life” Example of a Graph: $V =$ set of 6 people: Ali(1), Ahmed(2), Tarq(3), Youisf(4), Suha(5), and Saja(6), of ages 12, 15, 12, 15, 13, and 13, respectively.
- $E = \{ (x,y) \mid \text{if } x \text{ is younger than } y \}$

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Graph: Example_2

- $V = \{ A, B, C, D, E \}$
- $E = \{ (A,B), (B,C), (A,D), (D,A), (C,C), (E,D) \}$
- When (x,y) is an edge,
- we say that x is adjacent to y , and y is adjacent from x .
- A is adjacent to B .
- B is not adjacent to A .
- C is adjacent from B .

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The Graph Terminology

1. **Adjacent nodes:** two nodes are adjacent if they are connected by an edge

Ex: A is adjacent to B.



2. **Path:** a sequence of vertices that connect two or more nodes in a graph.

3. **Weighted graph:** a graph in which each edge carries a value



4. **Complete graph:** a graph in which every vertex is directly connected to every other vertex.

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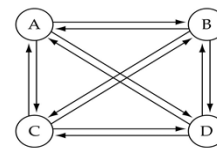
Graph Terminology

1. We can compute the number of edges in a complete directed graph with N vertices as follow:

$$\text{No. of } E = N * (N-1)$$

$$\text{No. of } E = 4 * (4-1)$$

$$\text{No. of } E = 12$$

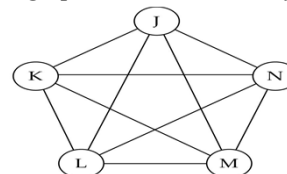


2. Also we can compute the number of edges in a complete undirected graph with N vertices as follow:

$$\text{No. of } E = N * (N-1) / 2$$

$$\text{No. of } E = 5 * (5-1) / 2$$

$$\text{No. of } E = 10$$



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Directed vs. Undirected Graphs

1. Directed graph (digraph)

- When the edges in a graph have a direction, the graph is called directed (or digraph).
- Let G be a directed graph
 1. The **Indegree** of a node x in G is the number of edges coming to x
 2. The **Outdegree** of x is the number of edges leaving x .

2. Undirected graph

- When the edges in a graph have no direction, the graph is called undirected

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Graph Representation

- For graphs to be computationally useful, they have to be conveniently represented in programs. There are two computer representations of graphs:
 1. **Adjacency matrix representation.**
 2. **Adjacency lists representation.**
- Basic Operations: Following are basic primary operations of a Graph:
 1. **Add Vertex:** Adds a vertex to the graph.
 2. **Add Edge:** Adds an edge between the two vertices of the graph.
 3. **Display Vertex:** Displays a vertex of the graph.

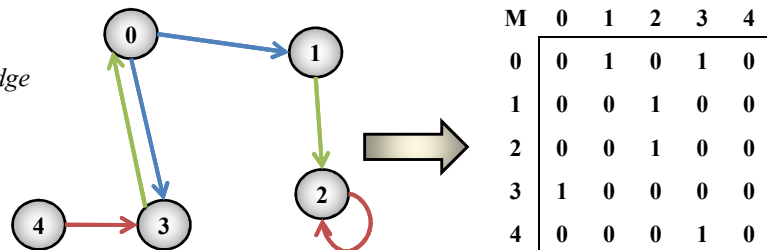
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Adjacency Matrix Representation

- In this representation, each graph of n nodes is represented by an $n \times n$ matrix A , that is, a two-dimensional array A :
- The nodes are (re)-labeled $0, 1, 2, \dots, n$
- $A[i][j] = 1$ if (i, j) is an edge
- $A[i][j] = 0$ if (i, j) is not an edge



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Adjacency Matrix Representation

- **Advantage:**
 - Simple to implement.
 - Easy and fast to tell if a pair (i, j) is an edge: simply check if $A[i][j]$ is 1 or 0
- **Disadvantage:**
 - No matter how few edges the graph has, the matrix takes $O(n^2)$ in memory.

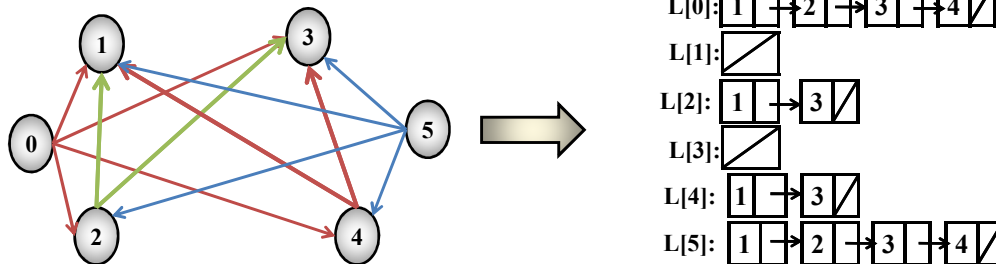
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Adjacency Lists Representation

- A graph of n nodes is represented by a one-dimensional array L of linked lists, where:
 - $L[i]$ is the linked list containing all the nodes adjacent from node i .
 - The nodes in the list $L[i]$ are in no particular order.



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Adjacency Lists Representation

- Advantage:**
 - Saves on space (memory): the representation takes as many memory words as there are nodes and edge.
- Disadvantage:**
 - It can take up to $O(n)$ time to determine if a pair of nodes (i, j) is an edge: one would have to search the linked list $L[i]$, which takes time proportional to the length of $L[i]$.

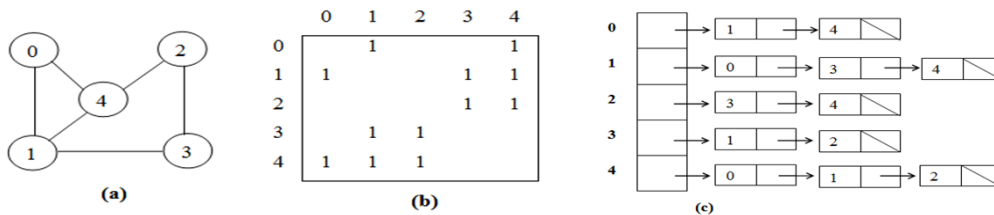
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Undirected Representation

- **Representations of the Undirected Graphs:** The same two representations for directed graphs can be used for undirected graphs
- **Adjacency matrix A:** $A[i][j]=1$ if (i,j) is an edge; 0 otherwise
- **Adjacency Lists:** $L[i]$ is the linked list containing all the neighbors of i

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Graph Traversal Techniques

- **There are two standard graph traversal techniques:**
 - **Depth-First Search (DFS)**
 - **Breadth-First Search (BFS)**
- In both **DFS** and **BFS**, the nodes of the undirected graph are visited in a systematic manner so that every node is visited exactly one.
- Both **BFS** and **DFS** give rise to a tree:
 - When a node x is visited, it is labeled as visited, and it is added to the tree.
 - If the traversal got to node x from node y , y is viewed as the parent of x , and x a child of y .

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Depth-First Search

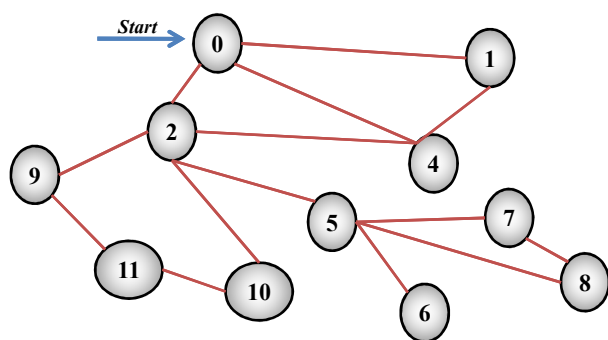
- *DFS follows the following rules:*
 1. Select an unvisited node x , visit it, and treat as the current node .
 2. Find an unvisited neighbor of the current node, visit it, and make it the new current node;
 3. If the current node has no unvisited neighbors, backtrack to the its parent, and make that parent the new current node;
 4. Repeat steps 3 and 4 until no more nodes can be visited.
 5. If there are still unvisited nodes, repeat from step 1.

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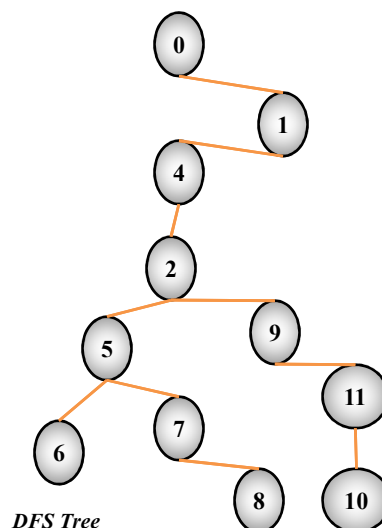


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Depth-First Search



Graph G



DFS Tree

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Breadth-First Search

- BFS follows the following rules:**

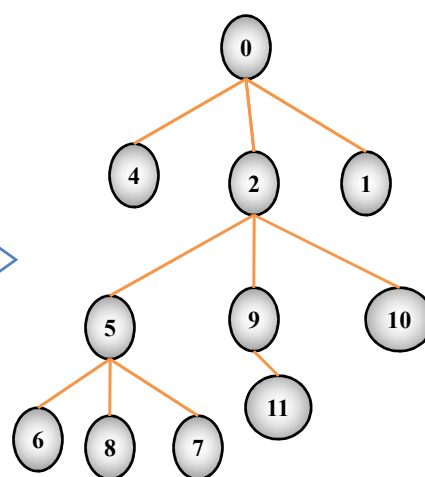
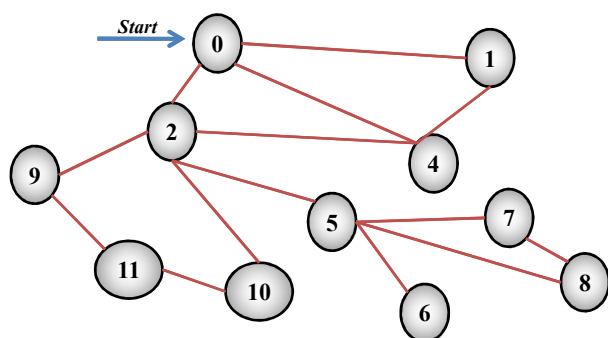
1. Select an unvisited node x , visit it, have it be the root in a BFS tree being formed. Its level is called the current level.
2. From each node z in the current level, in the order in which the level nodes were visited, visit all the unvisited neighbors of z . The newly visited nodes from this level form a new level that becomes the next current level.
3. Repeat step 2 until no more nodes can be visited.
4. If there are still unvisited nodes, repeat from Step 1.

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Breadth-First Search



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The End

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